

MSGSÜ, MAT 221, Sınav, 10 Kasım 2014, Saat 13:00, David Pierce. *Sadece 4 soruyu cevaplayın.*

As in class, \mathbb{N} is the set $\{1, 2, 3, \dots\}$ of **natural numbers**, and x' is the **successor** of x , so $1' = 2$, $2' = 3$, and so on. We also let $\omega = \{0\} \cup \mathbb{N}$ and $0' = 1$.

Problem 1. For a given value of n in \mathbb{N} , let \bar{x} denote the congruence-class of x modulo n , and let $\mathbb{Z}_n = \{\bar{x} : x \in \mathbb{N}\} = \{\bar{1}, \dots, \bar{n}\}$. If $\bar{x}' = \bar{y}'$, then $\bar{x} = \bar{y}$. Therefore we can define $\bar{x}' = \bar{x}'$. The structure $(\mathbb{Z}_n, \bar{1}, ')$ allows proofs by induction. We have shown that addition and multiplication on \mathbb{Z}_n can be defined recursively by

$$\begin{aligned} \bar{x} + \bar{1} &= \bar{x}', & \bar{x} + \bar{y}' &= (\bar{x} + \bar{y})', \\ \bar{x} \cdot \bar{1} &= \bar{x}, & \bar{x} \cdot \bar{y}' &= \bar{x} \cdot \bar{y} + \bar{x}. \end{aligned}$$

- (a) If $n = 6$, show that there is an operation on \mathbb{Z}_n given by

$$\bar{x}^{\bar{1}} = \bar{x}, \quad \bar{x}^{\bar{y}'} = \bar{x}^{\bar{y}} \cdot \bar{x}. \quad (*)$$

It is enough to fill out the table [in the solution].

- (b) If $n = 3$, show that there is *no* operation on \mathbb{Z}_n as in (*).

Solution.

$\bar{x}^{\bar{y}}$		y					
		1	2	3	4	5	6
(a)	x	1	1	1	1	1	1
		2	2	4	2	4	2
		3	3	3	3	3	3
		4	4	4	4	4	4
		5	5	2	5	2	5
		6	6	6	6	6	6

From the table it should be clear that, modulo 6,

$$a \equiv b \implies x^a \equiv x^b.$$

- (b) Using (*), we obtain $\bar{2}^{\bar{1}} = \bar{2}$, $\bar{2}^{\bar{2}} = \bar{4} = \bar{1}$, $\bar{2}^{\bar{3}} = \bar{2}$, so $\bar{2}^{\bar{4}} = \bar{1}$, so $\bar{2}^{\bar{1}} \neq \bar{2}^{\bar{4}}$, although $\bar{1} = \bar{4}$.

Problem 2. Using only the recursive definition of addition on \mathbb{N} and induction, prove that addition is associative.

Solution. We want to show $x + (y + z) = (x + y) + z$.

1. By the recursive definition of addition (namely $x + 1 = x'$ and $x + y' = (x + y)'$), we have

$$x + (y + 1) = x + y' = (x + y)' = (x + y) + 1,$$

so the claim is true when $z = 1$.

2. Suppose the claim is true when $z = t$. Then

$$\begin{aligned} x + (y + t') &= x + (y + t)' && \text{[def'n of +]} \\ &= (x + (y + t))' && \text{[def'n of +]} \\ &= ((x + y) + t)' && \text{[inductive hypothesis]} \\ &= (x + y) + t', && \text{[def'n of +]} \end{aligned}$$

so the claim is true when $z = t'$.

By induction, the claim holds for all z in \mathbb{N} (for all x and y in \mathbb{N}).

Problem 3. We know $2 \cdot \sum_{k=1}^n k = n \cdot (n + 1)$. For all n in \mathbb{N} , prove

$$\left(\sum_{k=1}^n k \right)^2 = \sum_{k=1}^n k^3.$$

Solution. The claim is obvious when $n = 1$. Suppose it is true when $n = m$. Then

$$\begin{aligned} \left(\sum_{k=1}^{m+1} k \right)^2 &= \left(\sum_{k=1}^m k \right)^2 + 2 \cdot \sum_{k=1}^m k \cdot (m + 1) + (m + 1)^2 \\ &= \sum_{k=1}^m k^3 + m \cdot (m + 1)^2 + (m + 1)^2 = \sum_{k=1}^m k^3 + (m + 1)^3 = \sum_{k=1}^{m+1} k^3. \end{aligned}$$

Problem 4. We can define the so-called binomial coefficients recursively by

$$\begin{aligned} \binom{0}{0} &= 1, & \binom{0}{k+1} &= 0, \\ \binom{n+1}{0} &= 1, & \binom{n+1}{k+1} &= \binom{n}{k} + \binom{n}{k+1}. \end{aligned}$$

Using only this definition, and induction, show that, for all n in ω ,

$$\sum_{k=0}^n \binom{n}{k} = 2^n.$$

Solution. The claim is obvious when $n = 0$. Suppose it is true when $n = m$. Then

$$\begin{aligned} \sum_{k=0}^{m+1} \binom{m+1}{k} &= \binom{m+1}{0} + \sum_{k=0}^m \binom{m+1}{k+1} \\ &= 1 + \sum_{k=0}^m \left(\binom{m}{k} + \binom{m}{k+1} \right) \\ &= \sum_{k=0}^m \binom{m}{k} + \sum_{k=0}^{m+1} \binom{m}{k} \\ &= \sum_{k=0}^m \binom{m}{k} + \sum_{k=0}^m \binom{m}{k} = 2 \cdot 2^m = 2^{m+1} \end{aligned}$$

(since $\binom{m+1}{0} = 1 = \binom{m}{0}$ and $\binom{m}{m+1} = 0$). Therefore, by induction, the claim is true for all n in ω . (We do not have $\binom{m}{m+1} = 0$ immediately from the definition, but can prove by induction that $\binom{m}{n} = 0$ when $n > m$.)

Problem 5. Let d be the greatest common divisor of 385 and 168.

(a) Find d .

(b) Find a solution from \mathbb{N} of one of the equations

$$385x = 168y + d, \quad 168x = 365y + d.$$

Solution.

(a) Since

$$385 = 168 \cdot 2 + 49$$

$$168 = 49 \cdot 3 + 21$$

$$49 = 21 \cdot 2 + 7$$

and $7 \mid 21$, we conclude $\gcd(385, 168) = 7$.

(b) From the previous computations,

$$\begin{aligned} 7 &= 49 - 21 \cdot 2 \\ &= 49 - (168 - 49 \cdot 3) \cdot 2 \\ &= 49 \cdot 7 - 168 \cdot 2 \\ &= (385 - 168 \cdot 2) \cdot 7 - 168 \cdot 2 \\ &= 385 \cdot 7 - 168 \cdot 16. \end{aligned}$$

Thus $385 \cdot 7 = 168 \cdot 16 + 7$.