The Hrushovski Programme

Alexandre Borovik (Unfinished) joint projects with Omaima Alshanqiti, Pınar Uğurlu, and Şükrü Yalçınkaya

Antalya Algebra Days XIV

16 May 2012

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Outline

The Steinberg Endomorphisms

Black Box Groups

Some model theory

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The Larsen-Pink Theorem

Groups with count function

The Hrushovski Programme

- The Steinberg Endomorphisms

Simple algebraic groups

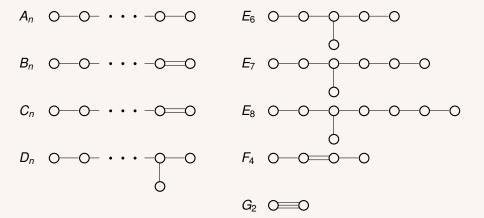
Chevalley:

A simple algebraic group is one of the following types:

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 A_n , B_n , C_n , D_n (classical groups) E_6 , E_7 , E_8 , F_4 , G_2 (exceptional groups)

Dynkin diagrams of simple algebraic groups Classical Groups Exceptional Groups



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The Steinberg Endomorphisms

- G simple algebraic group defined over \mathbb{F}_p
- σ rational endomorphism of *G* with finite group of fixed points

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 G_{σ} group of fixed points of σ

Example: Frobenius map induced by $x \mapsto x^q$, $q = p^k$.

Classification of Finite Simple Groups

Every non-abelian finite simple group is one of:

- 26 sporadic groups;
- alternating groups;
- *O*^{p'}(*G*_σ) (generated in *G*_σ by *p*-elements): groups of Lie type.

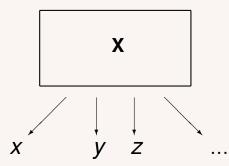
Uniform description of finite groups of Lie type

- ▶ for $T \sigma$ -invariant torus (Borel) in G form T_{σ} ,
- for $B \sigma$ -invariant Borel subgroup in G form B_{σ} , etc.

Lang-Steinberg: σ -invariant Borel subgroups do exist, etc.

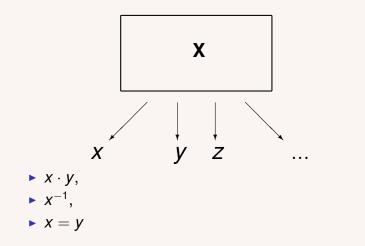
This is **THE** correct way to look at finite simple groups.

Black box groups



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Black box groups



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Matrix groups over finite fields

S a small set of invertible matrices over a finite field

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- $X = \langle S \rangle \leq \operatorname{GL}_n(q)$
- Input length: |S|n² log q

Matrix Groups

Let $X = \langle x_1, \ldots, x_n \rangle \leq GL_n(q)$ be a **big** matrix group so that |X| is astronomical.

Statistical study of random products of x₁,..., x_n is the only known approach to identification of X.

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- Determination of orders involves either
 - Factorization of integers into primes, or
 - Discrete logarithm problem over finite fields.

 Statistical study of 'random' products (Leedham-Green et al.) of

 x_1,\ldots,x_k

is the only known approach to identification of X.

- Basically, we are looking for a "short" and "easy to check by random testing" first order formula which identifies X.
- Existence /non-existence of elements of particular orders is an example.

Limits of crude statistical approach "Order of elements" approach fails for recognising

$$B_n(q) = \Omega_{2n+1}(q),$$

$$C_n(q) = PSp_{2n}(q),$$

q odd:

they have virtually the same statistics of orders of elements.

Here,

 $\Omega_{2n+1}(q)$ is the subgroup of index 2 in the orthogonal group $SO_{2n+1}(q)$,

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 $PSp_{2n}(q)$ is the projective symplectic group.

Why does statistics fail?

For large q, unipotent and non-semisimple elements occur with probability ~ 1/q and are "invisible": a random element is semisimple.

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Why does statistics fail?

Let $G = G(\overline{\mathbb{F}}_q)$ be a simple algebraic group.

- regular semisimple elements form an open subset of G
- statistics of orders of regular semisimple elements is determined by the **Dynkin diagram** of *G*, which is the same in the case of groups *B_n* and *C_n*, *n* ≥ 3:

$$BC_n, n \geq 2$$
 \longrightarrow

How one can fix the failure of statistics?

But the conjugacy classes and the structure of centralisers of *involutions* (elements of order 2) are determined by the **extended Dynkin diagrams** which are different:

$$\widetilde{B}_n, n \ge 3$$
 $\widetilde{C}_n, n \ge 3$ $\widetilde{C}_n, n \ge 3$

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How one can fix the failure of statistics?

(Extended) Dynkin diagrams are first order properties in the language of groups!

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Black-Box Curtis–Tits Theorem (Yalçinkaya)

Theorem

Let G be a (quasi)-simple black box group of (unknown) Lie type over a field of odd characteristic and known "global exponent" N: $g^N = 1$ for all $g \in G$.

There is a polynomial in log N algorithm which constructs the extended Dynkin diagram of G ...

... which also allows to construct "subgiagram" subgroups, etc.—in sort, to do a lot of fascinating stuff.

The moral of the story so far

Black box theory works much better ...

... if groups are studied up to elementary equivalence—rather than up to isomorphism

Some model theory

Elementary theory and elementary equiavalence

Let *G* be a group Th(G) the set of first order formulae true in *G*

Elementary equivalence:

$$G \equiv H \iff Th(G) = Th(H)$$

Pseudofinite groups

G is pseudofinite if

 every formula which is true on G is true on some finite group.

One may think of pseudofinite groups as ultraproducts of finite groups

$$G\simeq\prod_{i\in I}G_i/\mathcal{F}.$$

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Measure on G is the ultraproduct of canonical finite measures on G_i .

This is not a 0-1 measure!

There are sets of probability different from 0 and 1:

In PSL₂ over a field of odd order, formula

" $Z(C_G(x))$ contains an involution "

holds with probability $\approx 1/2$ (or 1/2 + infinitesimal).

Formulae like that make a decent approximation to the property

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"x has even order".

Some model theory

Uncountable categoricity

G is \aleph_1 -categorical $\iff \exists ! \ \widetilde{G} \equiv G$ of cardinality \aleph_1

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Some model theory

Definable set

Definable set: defined by a first order formula

$$C_G(a) = \{ x : ax = xa \},$$

 $a^G = \{ x : \exists y \, x = a^y \}.$

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Groups of finite Morley rank:

have a rank function

{ Definable sets in G^n } $\xrightarrow{rk} \mathbb{N} \cup \{0\}$

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- behaves like dimension of Zariski closed sets
- axiomatised by natural axioms

In the case of simple groups:

 \aleph_1 -categorical \iff of finite Morley rank

Some model theory

The Cherlin-Zilber Conjecture (c. 1980):

A simple infinite group of finite Morley rank is isomorphic as an abstract group to an algebraic group over an algebraically closed field.

The Hrushovski Programme

The Hrushovski Programme

The Hrushovski Programme *G* simple group of finite Morley rank ψ a generic automorphism Then $G_0 = C_G(\psi)$ is pseudofinite or at least behaves like pseudofinite.

In "real life", due to a theorem by Hrushovski: If *G* is algebraic over an a.c. field then

- ϕ is generalised Frobenius, and
- ► G₀ = C_G(φ) is the group of points of G over a pseudofinite field.

Pınar Uğurlu:

G simple group of finite Morley rank

 α automorphism of ${\it G}$

 $d(C_H(\alpha^{km})) = H$ for every connected α^k -invariant $H \leq G$ and every $k, m \in \mathbb{N}$.

 $C_G(\alpha^k)$ is pseudofinite for all $k \in \mathbb{N}$.

Then G is algebraic.

Proof does not use CFSG (the Classification of Finite Simple Groups).

- The Hrushovski Programme

Why CFSG has to be eliminated?

There is a good algebraic characterisation of pseudofinite fields:

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- perfect
- exactly one extension of every degree
- pseudo algebraically closed

but nothing of this kind is known for groups.

Larsen and Pink, 1998

For every *n* there exists a constant *J* depending only on *n* such that for any finite simple group *X* possessing a faithful linear or projective representation of dimension *n* over a field k we have either

(a)
$$|X| < J(n)$$
, or

(b) p := char(k) is positive and X is a group of Lie type in characteristic p.

- The Larsen-Pink Theorem

Larsen and Pink, equivalent statement:

A definably simple infinite pseudofinite subgroup $G \leq GL_n$ is a Chevalley group over a pseudofinite field.

- The Larsen-Pink Theorem

Proof in odd characteristic

• Work in the pair $G < \overline{G}$, where G is pseudofinite and \overline{G} is its Zariski closure (in GL_n).

- No use of CFSG.
- Use of large "definable" fragments of CFSG, for example:
 - Component analysis in groups of odd type.
 - Signalizer functor theory.

Count functions: motivation

- An attempt to replace both "finite" and "pseudofinite" by an unifying algebraic concept.
- We need to balance:
 - feasibility: the property needs to be verifiable in the context of the Hrushovski Programme
 - power: has to be strong enough to allow classification of definably simple groups with this property.

What follows is just a first try to achieve *power*; the feasibility was not even considered.

Count functions, after Krajíček and Scanlon

Let *A* be an algebraic structure and D the set of definable subsets in all A^n , n = 1, 2, ...

Let *R* be a linearly ordered unital commutative ring. A function

$$\mu: \mathcal{D} \to \boldsymbol{R}$$

is *a count function* on *A* over *R* if and only if it satisfies the following conditions.

Groups with count function

Count functions, continued

1.
$$\mu(\{a\}) = 1$$
 for any $a \in A^k$.

- 2. $\mu(X \cup Y) = \mu(X) + \mu(Y)$, whenever $X, Y \in D$ and X, Y are disjoint.
- 3. $\mu(X \times Y) = \mu(X) \times \mu(Y)$, whenever $X, Y, X \times Y \in \mathcal{D}$.
- 4. $\mu(X) = \mu(Y)$, whenever $X, Y \in \mathcal{D}$ and there is a definable bijection between X, Y.
- 5. $\mu(X) = c \cdot \mu(Y)$, whenever $c \in R, X, Y \in D$, and there is a definable map $f : X \longrightarrow Y$ such that each of its fibers $f^{(-1)}(y)$, where $y \in Y$, has count $\mu(f^{(-1)}(y)) = c$.

6.
$$\mu(X) \ge 0$$
 for all $X \in \mathcal{D}$.

A count function is *nontrivial* if 0 < 1 and the image of μ is not just $\{0\}$.

Tallied structures

For brevity, a structure with a nontrivial count function is called **tallied**.

Krajíček: Let A_i , for $i \in I$, be structures of the same languages, and assume that A is an ultraproduct of A_i . Assume that all A_i are tallied. Then A is tallied.

The Hrushovski Programme

Groups with count function

Tallied fields

A field *F* is *quasi-finite* if *F* is perfect and has precisely one extension of each degree (in a fixed algebraic closure \tilde{F}).

Scanlon: Any field admitting a non-trivial count function is quasi-finite.

Frobenious groups

A group *G* is called *Frobenius* if it contains a non-trivial proper subgroup *H* such that

$$H \cap H^g = 1$$
 for all $g \in G \setminus H$;

H is called a Frobenius complement of G. The set

$$\mathcal{K} = \left\{ \mathcal{G} \setminus \bigcup_{g \in \mathcal{G} \setminus \mathcal{H}} \mathcal{H}^g \right\} \cup \{1\}.$$

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is called the Frobenius kernel of G.

A version of the Frobenius Complement Theorem

B-Alshanqiti: Assume G is a tallied Frobenius group with a definable Frobenius complement H and the Frobenius kernel K. In addition, assume that H contains an involution. Then

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- ► *K* is a definable normal subgroup of *G*.
- ▶ *K* is an abelian group.
- ► *H* contains exactly one involution.

Counting arguments work!